Problem 1 (3 parts, 30 points)
Part A (12 points) Consider a DRAM chip organized as 4 million addresses of 64 -bit words. Assume both the DRAM cell and the DRAM chip are square. The column number and offset concatenate to form the memory address. Using the organization approach discussed in class, answer the following questions about the chip. Express all answers in decimal (not powers of two).

| number of columns | $4 \mathrm{M} \times 64=2^{22} \times 2^{6}=2^{28} ; 2^{14}=16 \mathrm{~K}$ |
| :---: | :---: |
| number of words per column | $2^{14} / 2^{6}=2^{8}=256$ |
| column decoder required ( $n$ to $m$ ) | 14 to 16K |
| total number of bits in address | $\log _{2}(4 M)=22$ |
| type of mux required ( $n$ to $m$ ) | 256 to 1 |
| number of address lines in column offset | 8 |

Part B (10 points) Consider a memory system with $\mathbf{1 6}$ million addresses of 32 -bit words using a 2 million address by 8 -bit word memory DRAM chip.

| word address lines for memory system | $\log _{2}(16 M)=24$ |
| :---: | :---: |
| chips needed in one bank | $32 / 8=4$ |
| banks for memory system | $16 M / 2 M=8$ |
| memory decoder required $(n$ to $m)$ | 3 to 8 |
| DRAM chips required | $4 \times 8=32$ |

Part C ( 8 points) Design a 32 million address by 8 bit memory system with four 16M x 4 memory chips. Label all busses and indicate bit width. Assume R/W is connected and not shown here. Use a bank decoder if necessary.


Problem 2 (3 parts, 24 points)

## Datapath Elements

Part A (6 points) Suppose the following inputs (in hexadecimal) are applied to the 32-bit barrel shifter used in the datapath. Determine the output (in hexadecimal). Assume the shift amount is drawn from the 16-bit immediate value.

| Shift Type | Shift Amount | Input Value | Output Value |
| :---: | :---: | :---: | :---: |
| logical | $0 x F F F 4$ | EB25ACE7 | 5 ACE7000 |
| arithmetic | $0 \times 0008$ | CAB15317 | FFCAB153 |
| rotate | $0 \times 000 \mathrm{C}$ | DE2F1B36 | B36DE2F1 |

Part B (8 points) For each bitwise logical function specification below, determine the LF code (in hexadecimal) to correctly program the logical unit.

| X Y | Out |  |
| :---: | :---: | :---: |
| 0 | 0 | $\mathrm{LF}_{0}$ |
| 1 | 0 | $\mathrm{LF}_{1}$ |
| 0 | 1 | $\mathrm{LF}_{2}$ |
| 1 | 1 | $\mathrm{LF}_{3}$ |


| logical function | LF |
| :---: | :---: |
| $\bar{Y}$ | 3 |
| $Y \cdot \bar{X}$ | 4 |
| $X+\bar{Y}$ | B |
| $X \cdot Y$ | 8 |

Part C (10 points) Given the following finite state diagram, fill in the state table below. The current state variable is $S$ and can be one of two states ( 0 or 1 ) and the next state variable is NS.


| $S$ | $A / \bar{B}$ | $N S$ | $M$ | $N$ |
| :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 |
| 1 | 1 | 1 | 0 | 0 |

Give the Boolean expression for computing NS in terms of the current state and the input.
$\mathrm{NS}=$
$\overline{S \oplus A}$ or $\bar{S} \cdot B+S \cdot A$

Problem 3 (3 parts, 26 points)
Microcode
Using the supplied datapath, write microcode fragments to accomplish the following procedures. Express all values, except memory addresses, in hexadecimal notation. Use ' $X$ ' when a value is don't cared. For maximum credit, complete the description field. $\oplus$ means bitwise logical XOR. In each part, modify only registers 7 \& 8.
Part A (6 points) $\quad R_{7}=15 \cdot R_{8}$

| \# | $X$ | $Y$ | z | rwe | ${ }_{\text {im }}^{\substack{\text { en }}}$ | im va | au en | -a / s | lu en |  |  | st en | ${ }_{\text {r }}$ - | msel | description |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 8 | X |  | 1 | 1 | FFFC | 0 | X | 0 |  |  | 0 | X | 0 | R7 = R8 sll 4 |
| 2 | 7 | 8 | 7 | 1 | 0 | X | 1 | 1 | 0 |  |  | 0 | $x$ | 0 | R7 = R7-R8 |
| 3 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

Part B (12 points) Compute mem[3000] + 20 and store the result in mem[4000].

| $\#$ | $X$ | $Y$ | Z | rwe | im <br> $e n$ | im va | $a u$ <br> $e n$ | ca <br> s | lu <br> $e n$ | lf | su <br> $e n$ | st <br> $e n$ <br> $e n$ | st <br> $e n$ | $r /$ <br> $-w$ | msel | description |  |
| :--- | :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | x | x | 7 | 1 | 1 | 3000 | 0 | x | 1 | C | 0 | x | 0 | 0 | x | 0 | $\mathrm{R} 7=3000$ |
| 2 | 7 | x | 7 | 1 | 0 | x | 0 | x | 0 | x | 0 | x | 1 | 0 | 1 | 1 | $\mathrm{R} 7=\mathrm{Mem}[3000]$ |
| 3 | 7 | x | 7 | 1 | 1 | 14 | 1 | 0 | 0 | x | 0 | x | 0 | 0 | x | 0 | $\mathrm{R} 7=\mathrm{R} 7+20$ |
| 4 | x | x | 8 | 1 | 1 | 4000 | 0 | x | 1 | C | 0 | x | 0 | 0 | x | 0 | $\mathrm{R} 8=4000$ |
| 5 | 8 | 7 | x | 0 | 0 | x | 0 | x | 0 | x | 0 | x | 0 | 1 | 0 | 1 | $\mathrm{Mem}[4000]=\mathrm{R} 7$ |
| 6 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

Part C (8 points) Register 7 holds two packed 16 bit binary strings A and B as illustrated below.

|  |  |  | B |
| :--- | :--- | :--- | :--- | :--- |
| 31 | 16 | 15 | 0 |

Write a microcode sequence that unpacks A and B and computes $R_{8}=A \oplus B$.

| \# | $X$ | $Y$ | Z | rwe | $\begin{aligned} & \text { im } \\ & \text { en } \end{aligned}$ | im va | $\begin{aligned} & a u \\ & \text { en } \end{aligned}$ | $\begin{aligned} & -a \\ & 1 / 2 \end{aligned}$ | $\begin{gathered} \text { lu } \\ \text { en } \end{gathered}$ | If | su | st | $\begin{aligned} & \text { ld } \\ & \text { en } \end{aligned}$ | $\begin{aligned} & \text { st } \\ & e n \end{aligned}$ | r/ | msel | description |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 7 | x | 8 | 1 | 1 | FFFF | 0 | X | 1 | 8 | 0 | $x$ | 0 | 0 | x | 0 | $\mathrm{R} 8=\mathrm{R} 7$ and FFFF |
| 2 | 7 | x | 7 | 1 | 1 | 10 | 0 | x | 0 | x | 1 | 0 | 0 | 0 | x | 0 | $\mathrm{R} 7=\mathrm{R} 7 \mathrm{srl} 16$ |
| 3 | 7 | 8 | 7 | 1 | 0 | x | 0 | x | 1 | 6 | 0 | x | 0 | 0 | x | 0 | $\mathrm{R} 8=\mathrm{R} 7$ xor R 8 |
| 4 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

Assuming A and B are nonzero, what must be true about A and B for the result in $\mathrm{R}_{8}$ to be zero?

$$
A=B
$$

Part A (10 points) Design a toggle cell using two transparent latches, two 2 to 1 muxes, and one inverter. Your toggle cell should have an active high toggle enable input TE, and an active low clear input $\overline{C L R}$, clock inputs $\Phi_{1}$ and $\Phi_{2}$, and an output Out. The $\overline{C L R}$ signal has precedence over TE. Label all signals. Also complete the behavior table for the toggle cell.


Part B (10 points) Now combine these toggle cells to build a divide by seven counter. Your counter should have an external clear, external count enable, and three count outputs $\mathrm{O}_{2}, \mathrm{O}_{1}, \mathrm{O}_{0}$. Use any basic gates (AND, OR, NAND, NOR, XOR \& NOT) you require. Assume clock inputs to the toggle cells are already connected. Your design must support multi-digit systems.


